

## Problem C. Intermediate Verticality

Time limit: 1 second  
 Memory limit: 512 megabytes

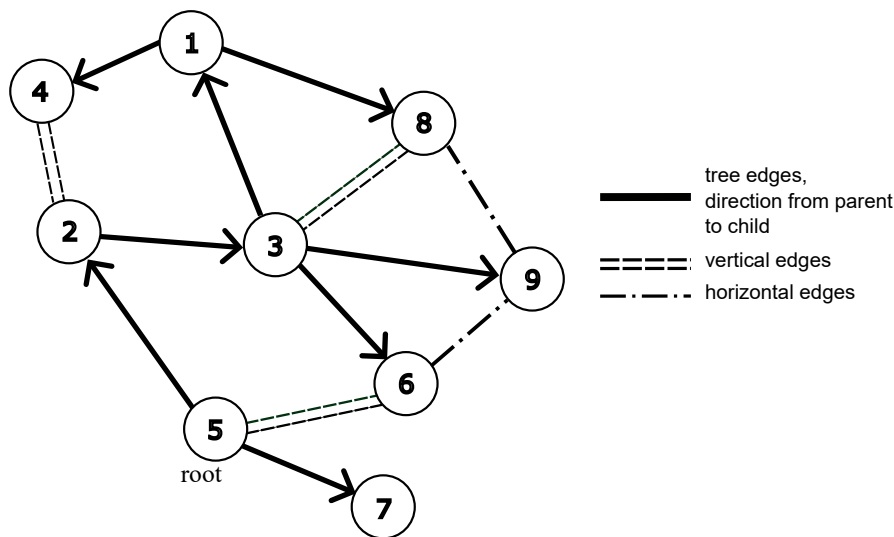
Two classical graph algorithms — depth-first search and breadth-first search — construct two spanning trees in a graph. The depth-first search is known for producing a tree that has no *horizontal* edges, which are edges connecting vertices that are not ancestors of each other, while breadth-first search is known for producing a tree that has no *vertical* edges — edges connecting a vertex to its ancestor in the tree. In this problem, you will need to construct an intermediate spanning tree that has a specified number of horizontal and vertical edges.

Recall that an undirected graph consists of a set of vertices  $V$  and a set of edges  $E$ , where each edge connects two vertices. We will consider connected graphs, where it is possible to reach any vertex from any other vertex via edges. A tree is a connected undirected graph that contains no cycles, and a spanning tree in a graph is a subset of its edges that forms a tree, allowing one to reach any vertex of the graph from any other. Recall two fundamental properties of a tree: in a tree with  $n$  vertices, there are exactly  $n - 1$  edges, and there is exactly one path between any two vertices in a tree.

We will designate a vertex  $r$  in the graph, which we will call the *root* of the tree. The vertices that lie on the unique path from vertex  $x$  to vertex  $r$  are called the *ancestors* of vertex  $x$ , and the first vertex on this path is called the *parent* of vertex  $x$  and is denoted as  $p_x$ . The root has no parent.

If a root and a spanning tree are fixed in the graph, then all edges of the graph can be divided into three types:

- *tree edges* — the edges of the chosen spanning tree;
- *vertical edges* — the edges not belonging to the tree that connect a vertex to its ancestor;
- *horizontal edges* — the remaining edges of the graph.



The *verticality* of the spanning tree in the graph is defined as the number of vertical edges.

You are given a graph with  $n$  vertices and  $m$  edges, the root of the tree  $r$ , and a number  $h$ ,  $0 \leq h \leq m - n + 1$ . You need to construct a spanning tree of the given graph with root at  $r$ , having a verticality equal to  $h$ , or report that such a tree does not exist.

### Input

Each test consists of several sets of input data. The first line contains one integer  $t$  — the number of sets of input data ( $1 \leq t \leq 10^5$ ). Following this are descriptions of  $t$  sets of input data.

In the first line of each set of input data, there are four integers  $n$ ,  $m$ ,  $r$ , and  $h$  — the number of vertices and edges in the graph, respectively, the index of the root vertex, and the required verticality of the future spanning tree ( $2 \leq n \leq 3 \cdot 10^5$ ;  $n - 1 \leq m \leq 3 \cdot 10^5$ ;  $1 \leq r \leq n$ ;  $0 \leq h \leq m - n + 1$ ).

In each of the following  $m$  lines, there are two integers  $u_i, v_i$  — the indices of the vertices connected by an edge in the graph ( $1 \leq u_i, v_i \leq n$ ;  $u_i \neq v_i$ ).

It is guaranteed that all graphs are connected, contain no loops or multiple edges. It is guaranteed that the sum of  $n$  across all input data sets does not exceed  $3 \cdot 10^5$ . It is guaranteed that the sum of  $m$  across all input data sets does not exceed  $3 \cdot 10^5$ .

## Output

For each test case, find the required spanning tree  $T$  and output in a separate line  $n$  integers  $p_1, p_2, \dots, p_n$ , where  $p_i$  is the index of the parent of the  $i$ -th vertex in the tree  $T$  ( $1 \leq p_i \leq n$ ). For  $p_r$ , you can output any number from 1 to  $n$ . If a tree  $T$  with the desired properties does not exist, output  $n$  numbers  $-1$  instead.

## Example

standard input	standard output
4	2 1 2 3 3
5 7 2 0	2 1 4 1 1
1 2	2 1 5 5 1
2 3	2 1 4 1 3
3 4	
4 1	
3 5	
5 4	
1 5	
5 7 2 1	
1 2	
2 3	
3 4	
4 1	
3 5	
5 4	
1 5	
5 7 2 2	
1 2	
2 3	
3 4	
4 1	
3 5	
5 4	
1 5	
5 7 2 3	
1 2	
2 3	
3 4	
4 1	
3 5	
5 4	
1 5	